

(Early) Memory Corruption Attacks

CS-576 Systems Security

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Fall 2018

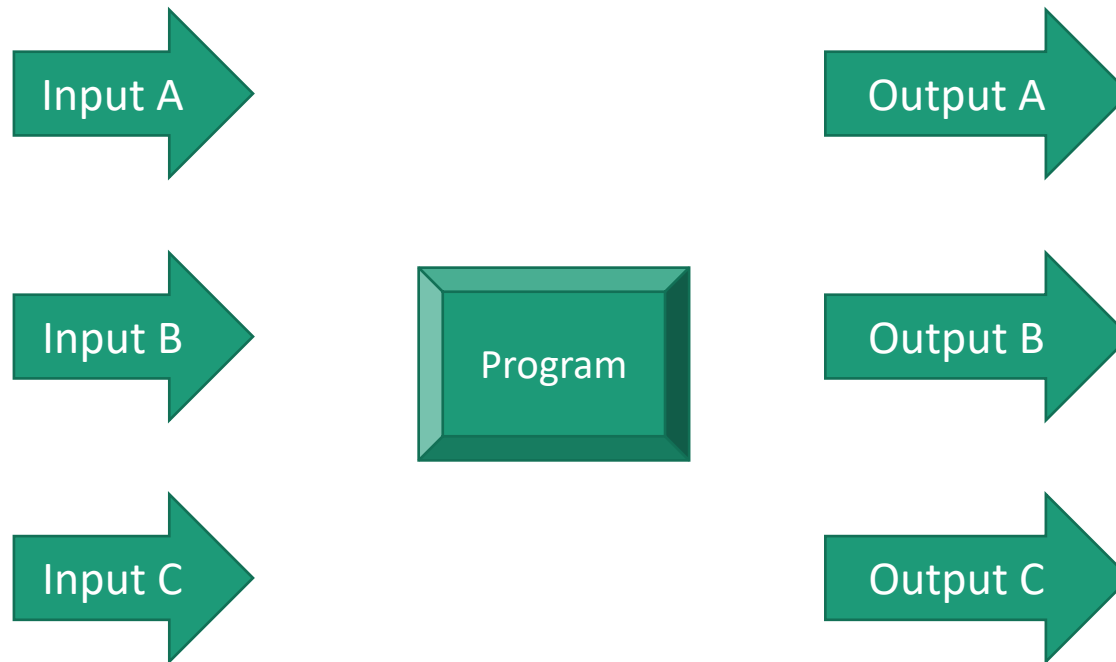
Memory Corruption

“Memory corruption occurs in a computer program when the contents of a memory location are unintentionally modified due to programming errors; this is termed **violating memory safety**.

When the corrupted memory contents are used later in that program, it leads either to program crash or to **strange and bizarre program behavior**. “

--wikipedia

Programs Are Deterministic

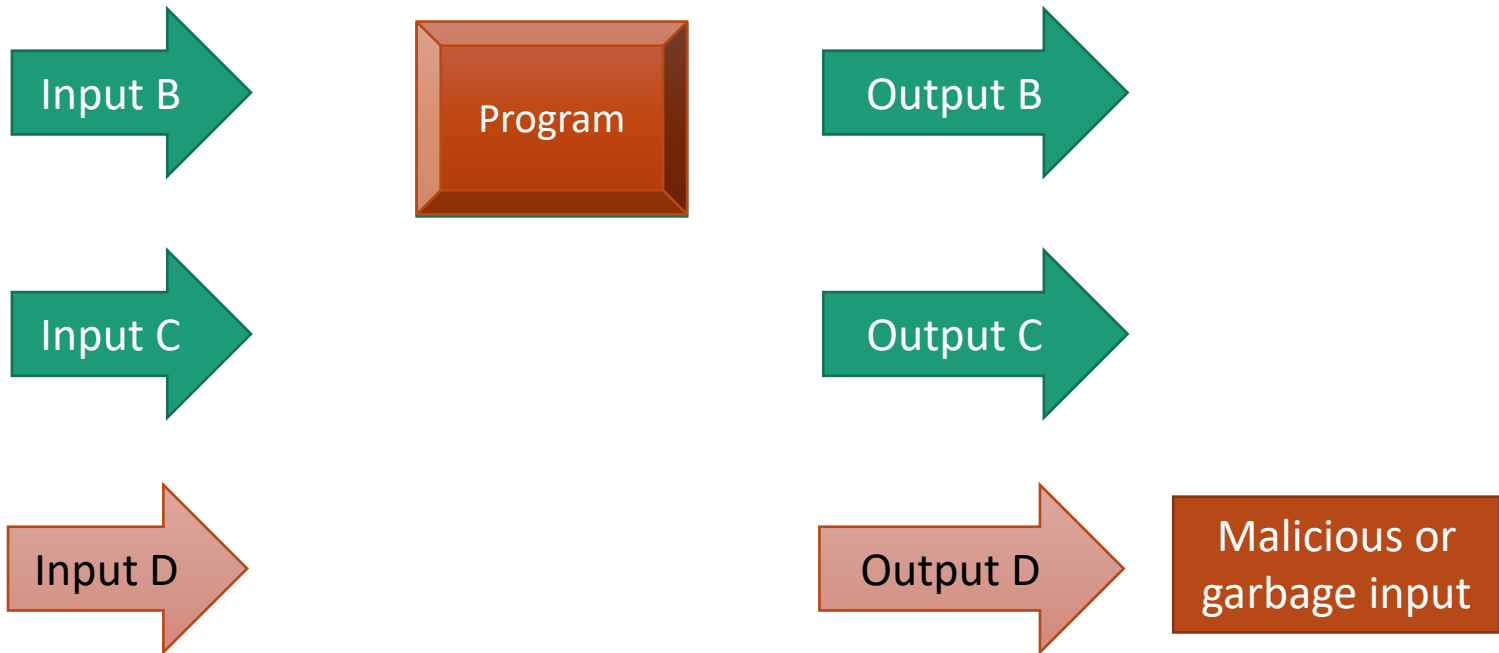


The program implements the functionality intended by the developer

Programs Are Deterministic

Program functionality severely altered

Original program → arbitrary program based on input



Unexpected or untested input triggering vulnerability in program

Malicious or garbage input

**Incorrect handling of
untested or incorrect input is
one the main causes of
software **vulnerabilities****

Remote Vs. local

Local attacks

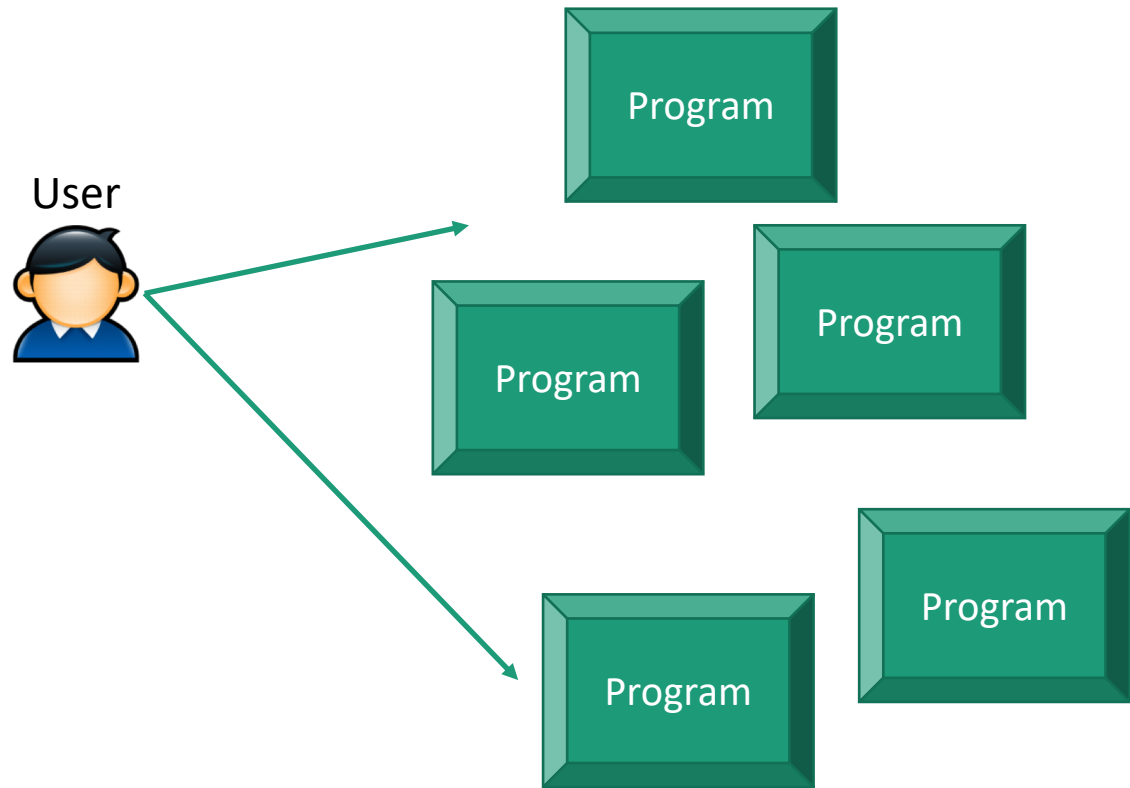
- If the user input can be only provided by a local user

Remote attack

- If the user input can be only provided over the network

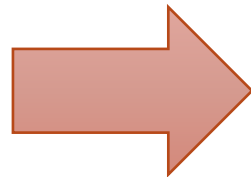
Local Attacks

All programs run with the privileges of the running user (Effective UID)

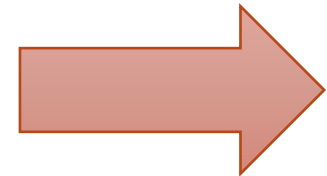


Local Attacks

Input produced by
another user



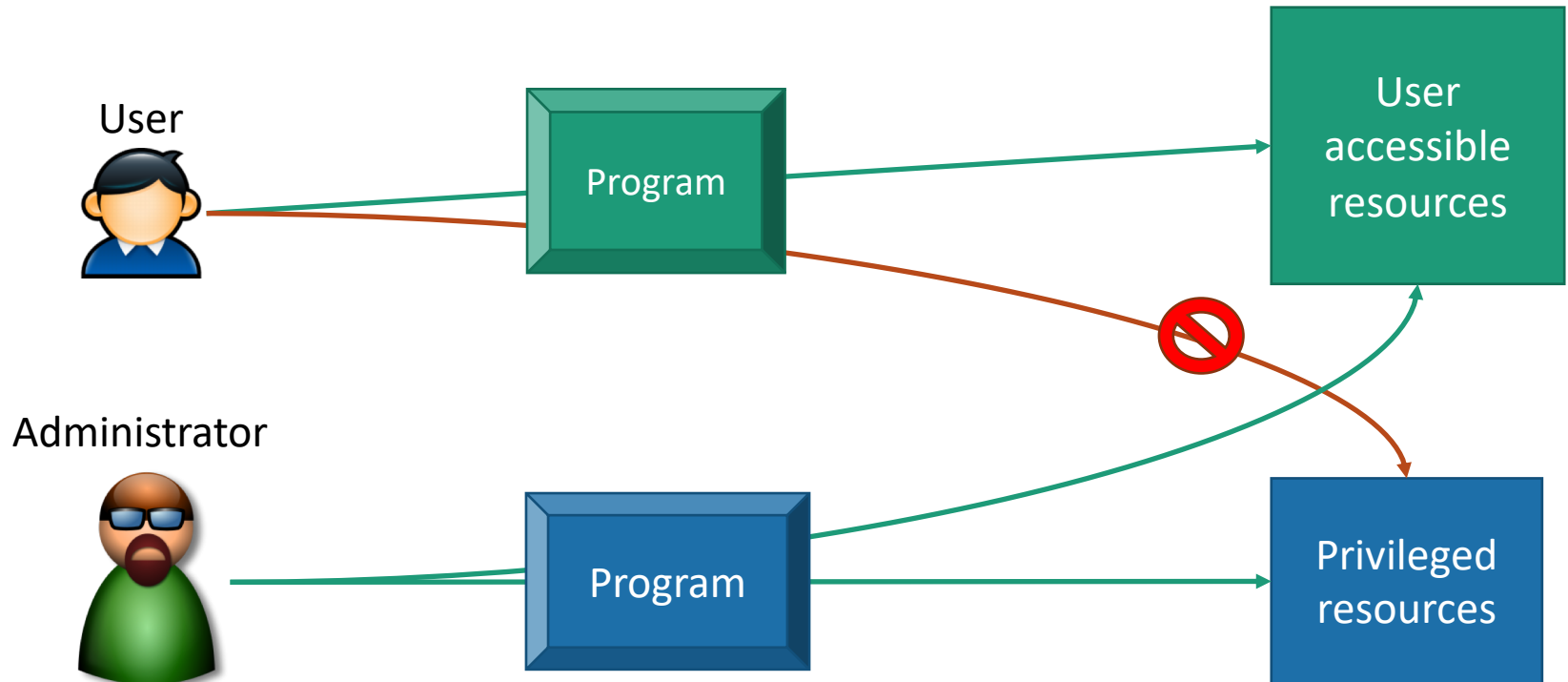
Program



Arbitrary program
executes with the rights
of the user executing it

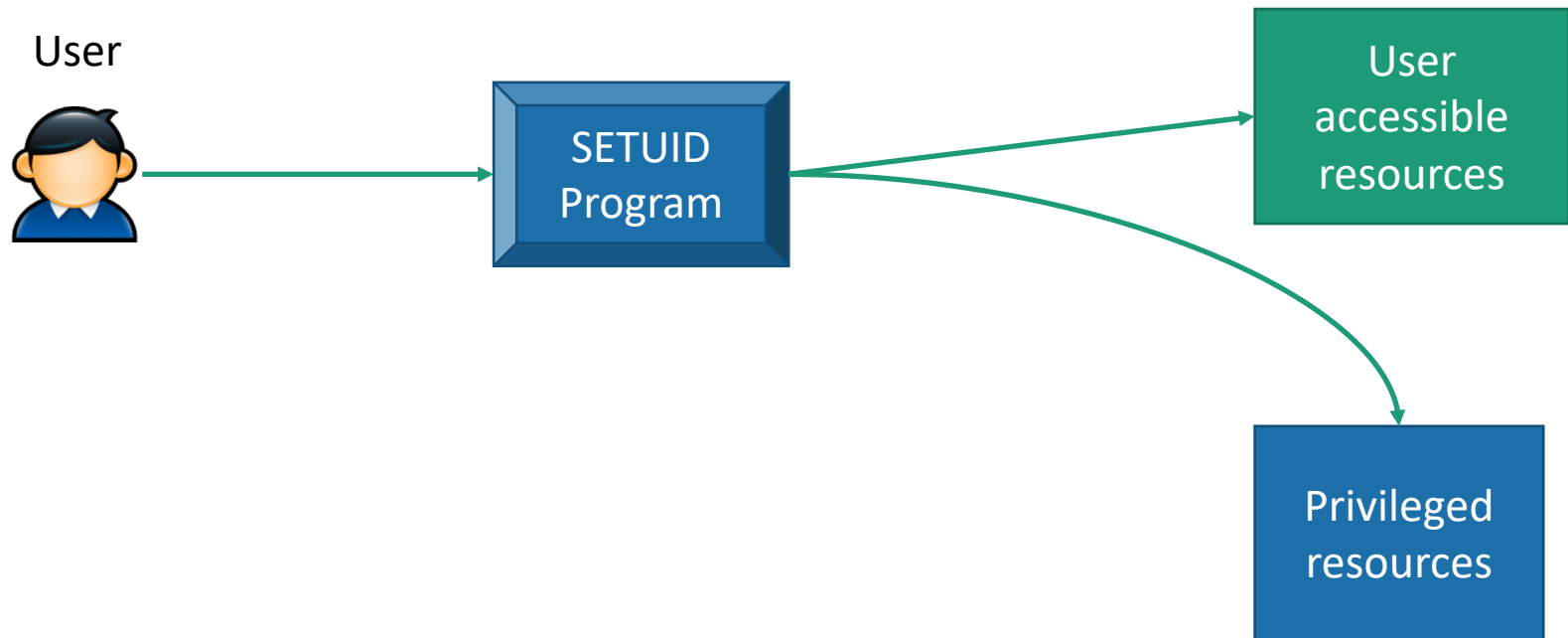
Privileged Resources

KERNEL



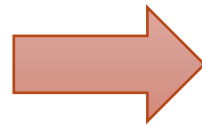
SETUID Programs

Programs that run with the privileges of their owner, not the executing user

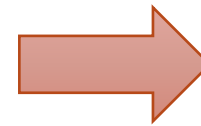


Privilege Escalation Attacks

Input produced by
another user



SETUID
Program



Privileged
resources

The arbitrary program
executes with elevated
privileges

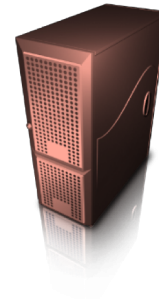
Remote Attacks

An arbitrary program is
run at a remote host

Bad Input

www.stevens.edu

Host: www
OS: Debian
HTTP Server: nginx



Remote Attacks

An arbitrary program is run at a remote host

Bad Input

With high privileges

www.stevens.edu

Host: www
OS: Debian
HTTP Server: nginxd



Running as root

Common Vulnerabilities

Overflows: Writing beyond the end of a buffer

Underflows: Writing beyond the beginning of a buffer

Use-after-free: Using memory after it has been freed

Uninitialized memory: Using pointer before initialization

Null pointer dereferences: Using NULL pointers

Type confusion: Assume a variable/object has the wrong type

HW errors: Hammering memory to cause bit flips to non-owned memory

CWE™ is a community-developed list of common software security weaknesses. It serves as a common language, a measuring stick for software security tools, and as a baseline for weakness identification, mitigation, and prevention efforts.

View the CWE List

[View by Research Concepts](#)

[View by Development Concepts](#)

[View by Architectural Concepts](#)

Search CWE

Easily find a specific software weakness by performing a search of the CWE List by keywords(s) or by CWE-ID Number. To search by multiple keywords, separate each by a space.

See the full [CWE List](#) page for enhanced information, downloads, and more.

Buffer Overflows

Buffer Overflows

Writing outside the boundaries of a buffer

Common programmer errors that lead to it ...

- Insufficient input checks/wrong assumptions about input
- Unchecked buffer size
- Integer overflows



Stack Overflows

Stack Overflow Example

```
int mytest(char *str)
{
    char buf[16];

    strcpy(buf, str);

    printf("%s\n", buf);

    return strlen(buf);
}
```

Stack Overflow Example

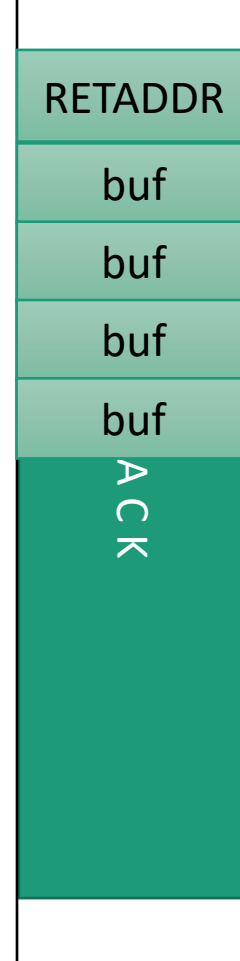
```
int mytest(char *str)
{
    char buf[16];

    strcpy(buf, str);

    printf("%s\n", buf);

    return strlen(buf);
}
```

High address/stack bottom



Low address/stack top

Stack Overflow Example

```
int mytest(char *str)
{
    char buf[16];

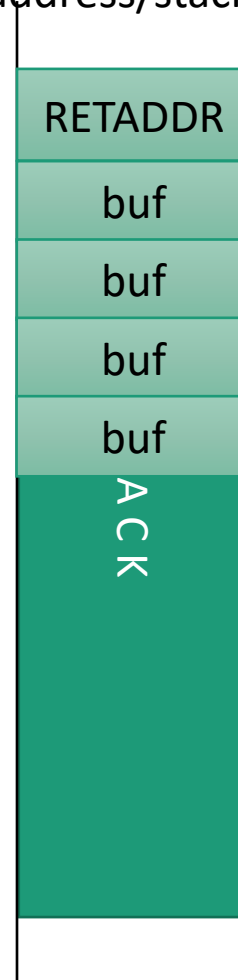
    strcpy(buf, str);

    printf("%s\n", buf);

    return strlen(buf);
}
```

```
./mytest AAAAA
```

High address/stack bottom



Low address/stack top

Stack Overflow Example

```
int mytest(char *str)
{
    char buf[16];

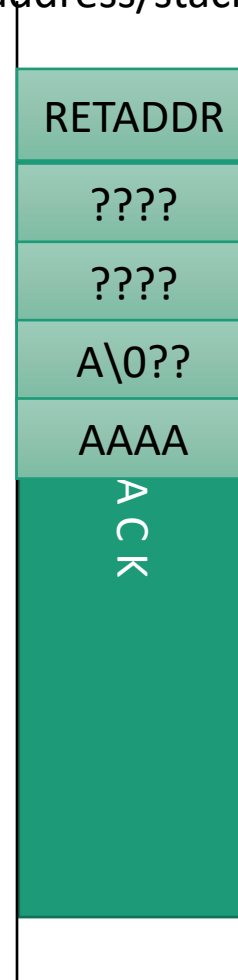
    strcpy(buf, str);

    printf("%s\n", buf);

    return strlen(buf);
}
```

```
./mytest AAAAA
```

High address/stack bottom



Low address/stack top

Stack Overflow Example

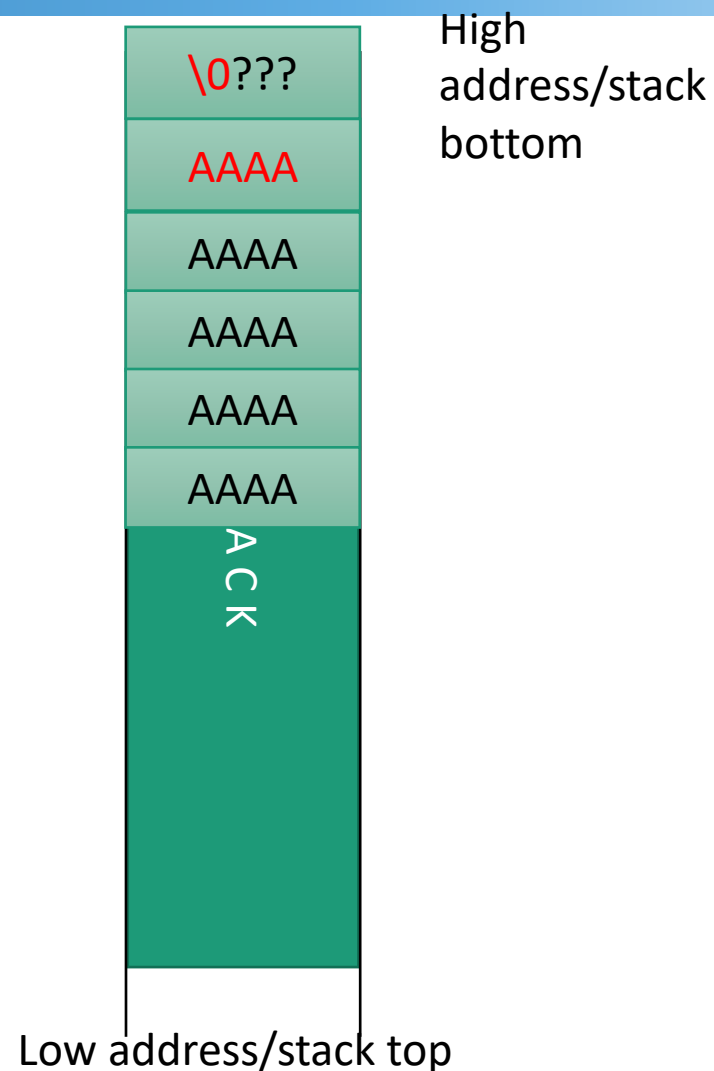
```
int mytest(char *str)
{
    char buf[16];

    strcpy(buf, str);

    printf("%s\n", buf);

    return strlen(buf);
}
```

```
./mytest AAAAAAAAAAAAAAAAAAAAAA
```



Stack Overflow Example

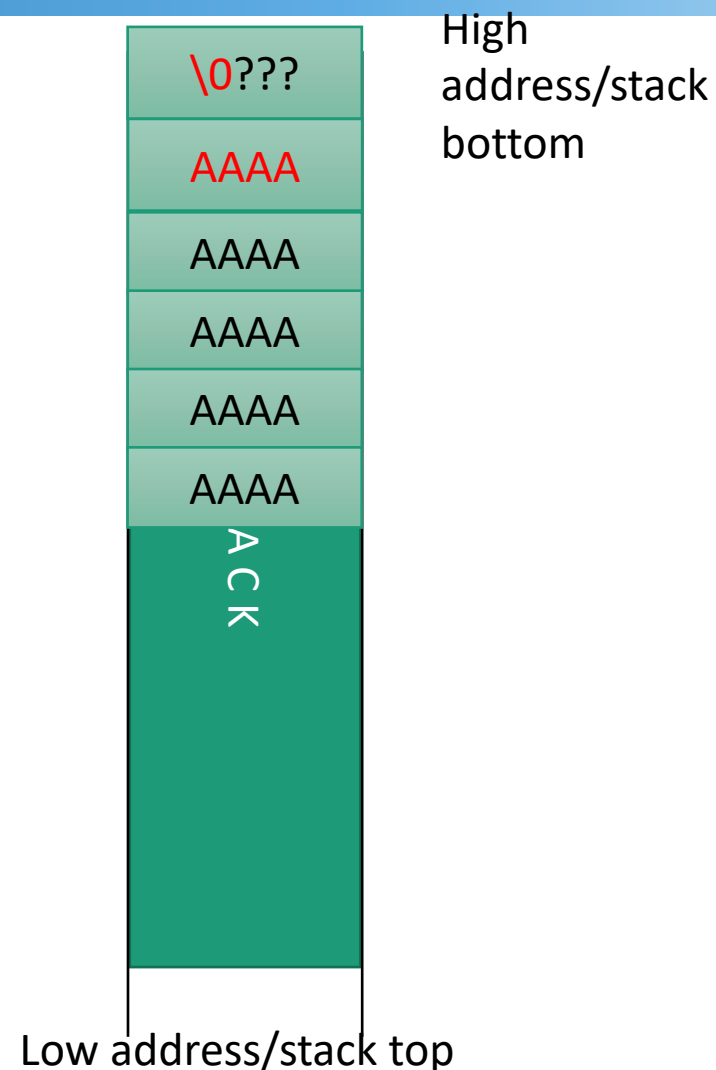
```
int mytest(char *str)
{
    char buf[16];

    strcpy(buf, str);

    printf("%s\n", buf);

    return strlen(buf);
}
```

```
./mytest AAAAAAAAAAAAAAAAAAAAAA
```



Control-Flow Hijacking

The saved return address is a code pointer stored in memory

- Controlling it grants control of a control-flow instruction (e.g., ret)

Untrusted inputs that lead to corruption of a code pointer lead to **control-flow hijacking attacks**

Other Code Pointers

Return
address

```
return;
```

```
ret
```

Function
address

```
typedef void (*cmpf_t)(int, int);  
void compare(int array[], int len, int num, cmpf_t f)  
{  
    int i;  
    for (i = 0; i < len; i++)  
        if (array[i] < num)  
            f(i, array[i]);  
}
```

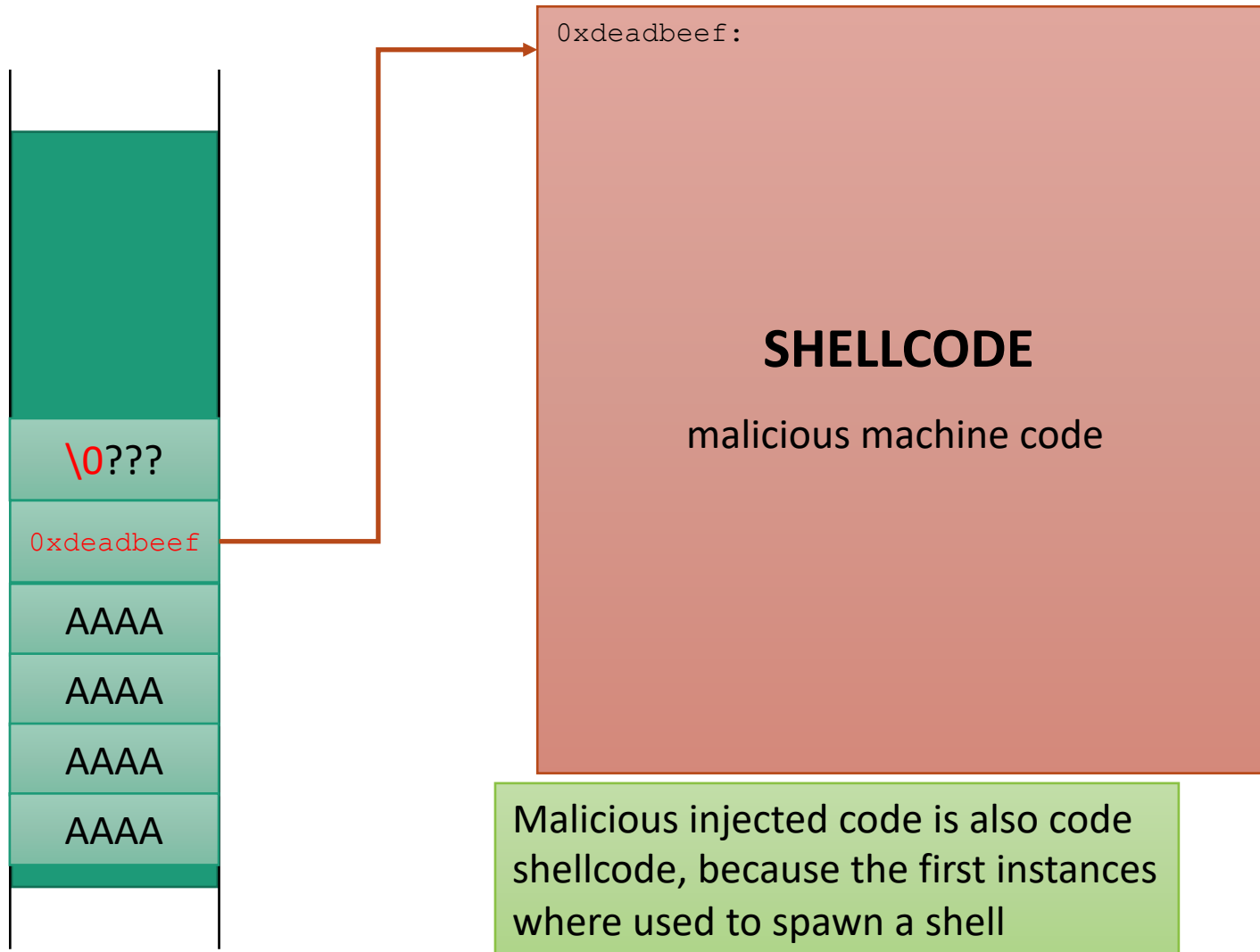
```
call *(rax)
```

Jump
table

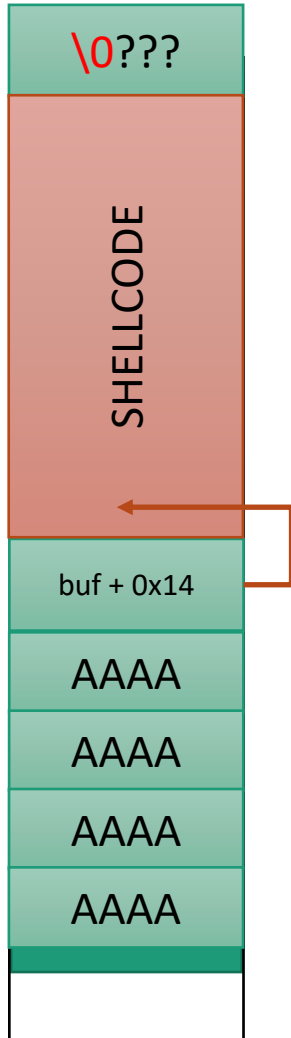
```
switch (option) {  
case 0:  
    Code ...  
case 1:  
    Code ...  
...  
}
```

```
jmp *(rax)
```

Where to Point Execution



Injecting Shellcode

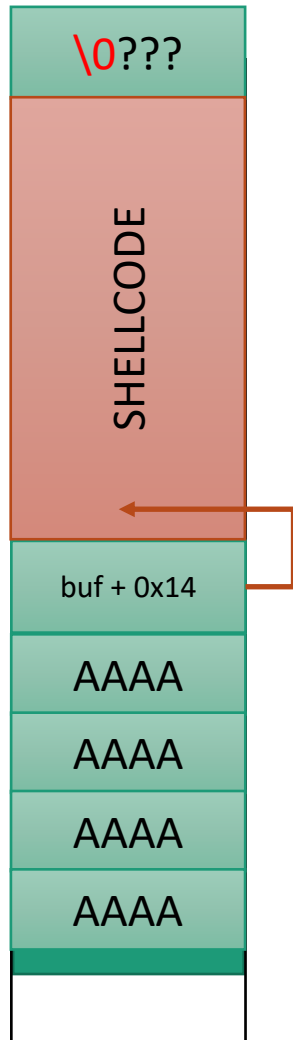


Code Injection

Code injection (CI) - Injecting machine code into a vulnerable program's memory

Code injections attacks inject code and use control-flow hijacking to execute that code

Shellcode Limitations



Injected shellcode cannot include a null byte because of `strcpy()`

Shellcode needs to be carefully crafted to avoid disallowed bytes

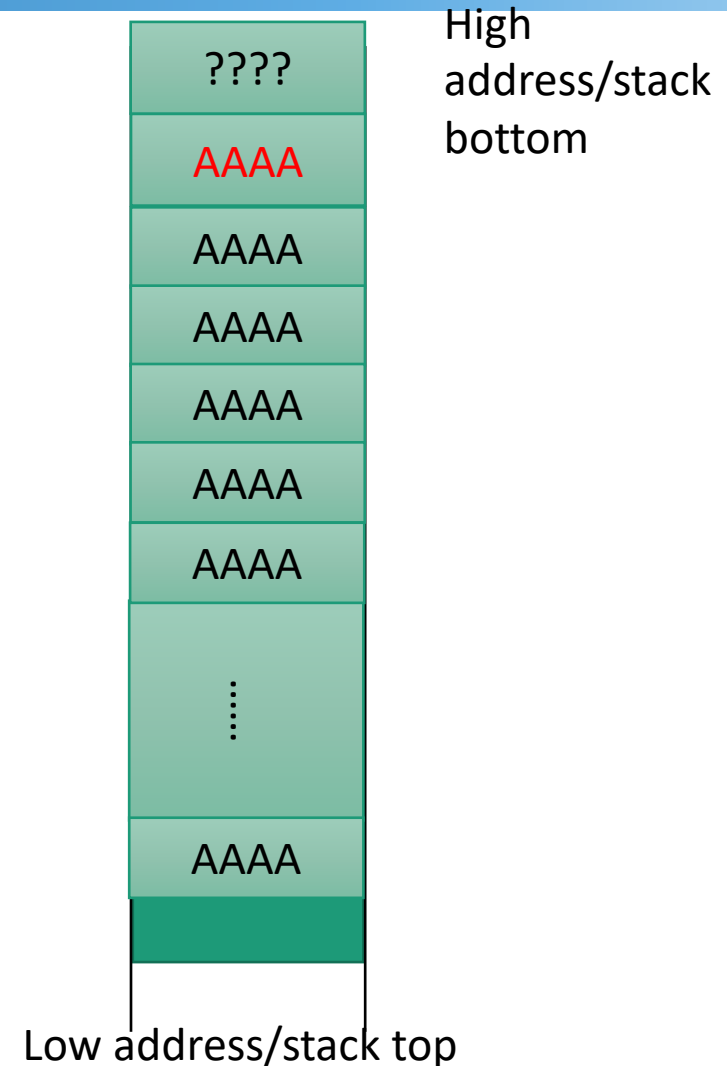
Other methods of copying data may not have the same limitation: `memcpy()`, `gets()`, `read()`, `fread()`, custom copy routines, etc.

Stack Overflow Using read()

```
static void getURL(void)
{
    char buf[64];

    read(STDIN_FILENO, buf, 128);
    get_webpage(buf);
}
```

No limitation on bytes read.



Stack Overflow with FP

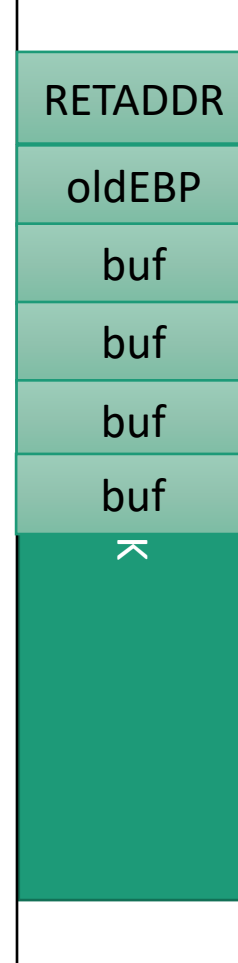
```
int mytest(char *str)
{
    char buf[16];

    strcpy(buf, str);

    printf("%s\n", buf);

    return strlen(buf);
}
```

High address/stack bottom



Low address/stack top

Stack Overflow with FP

```
int mytest(char *str)
{
    char buf[16];

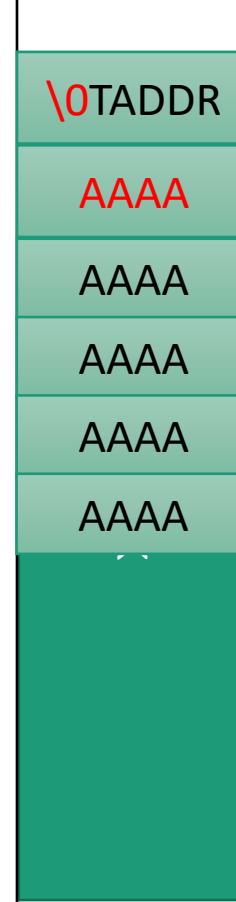
    strcpy(buf, str);

    printf("%s\n", buf);

    return strlen(buf);
}

./mytest AAAAAAAAAAAAAAAAAAAAAA
```

High address/stack bottom



Low address/stack top

Stack Overflow with FP

```
int mytest(char *str)
{
    char buf[16];

    strcpy(buf, str);

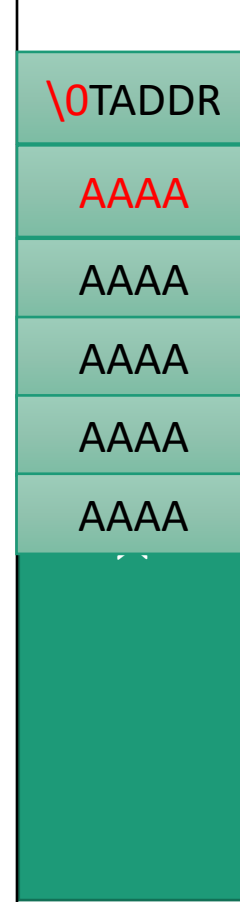
    printf("%s\n", buf);

    return strlen(buf);
}
```

```
./mytest AAAAAAAAAAAAAAAAAAAAAA
```

```
80484e1: c9          leave
80484e2: c3          ret
```

High address/stack bottom



Low address/stack top

Stack Overflow with FP

```
int mytest(char *str)
{
    char buf[16];

    strcpy(buf, str);

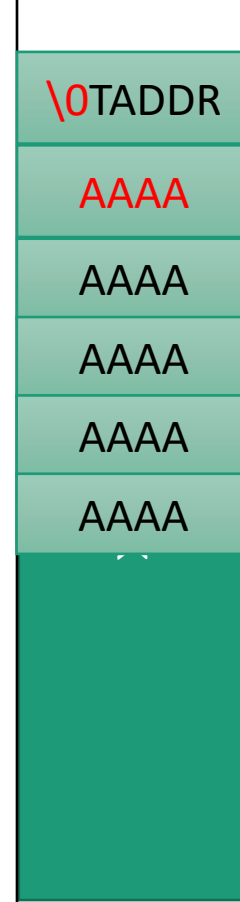
    printf("%s\n", buf);

    return strlen(buf);
}
```

```
./mytest AAAAAAAAAAAAAAAAAAAAAA
```

```
80480000 re
80480000
movl    %ebp, %esp
pop     %ebp
Function exit (LEAVE)
```

High address/stack bottom



Low address/stack top

Stack Overflow with FP

```
int mytest(char *str)
{
    char buf[16];

    strcpy(buf, str);

    printf("%s\n", buf);

    return strlen(buf);
}
```

```
./mytest AAAAAAAAAAAAAA\x3c\xca\xff\xffAAAA
```



High address/stack bottom

Low address/stack top

```
Function exit (LEAVE)
8048048: movl    %ebp, %esp
8048048: pop    %ebp
```